

# On the Stability Constant for $\varepsilon$ -light Vertex Subsets

## Upper Bound and Structural Reduction

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### Abstract

For a graph  $G = (V, E)$  with Laplacian  $L$ , a vertex subset  $S \subseteq V$  is  $\varepsilon$ -light if the spectral inequality  $\varepsilon L - L_S \succeq 0$  holds. We investigate the existence of a universal constant  $c > 0$  such that every graph contains an  $\varepsilon$ -light subset of size  $|S| \geq c\varepsilon|V|$ . We prove rigorously that the optimal constant satisfies  $c \leq 1/2$  via a perfect matching obstruction. Furthermore, we reduce the lower bound problem to a specific spectral subset selection hypothesis. We demonstrate that this hypothesis fails for general positive semidefinite matrices, thereby isolating the specific structural properties of graph Laplacians required to establish  $c = 1/2$ .

## 1 Introduction

Let  $G = (V, E)$  be a finite simple graph with  $n = |V|$  vertices. Let  $L$  be the Laplacian matrix of  $G$ . For any subset  $S \subseteq V$ , let  $L_S$  denote the Laplacian of the induced subgraph  $G_S = (V, E(S, S))$ .

**Definition 1** ( $\varepsilon$ -light subset). *A subset  $S \subseteq V$  is called  $\varepsilon$ -light if*

$$L_S \preceq \varepsilon L,$$

where  $\preceq$  denotes the Loewner order (i.e.,  $\varepsilon L - L_S$  is positive semidefinite).

Experimental evidence suggests that  $c = 1/2$  is the sharp constant. This note establishes the upper bound  $c \leq 1/2$  and provides a rigorous reduction of the lower bound to a matrix discrepancy conjecture on the image of the Laplacian.

## 2 Upper Bound: $c \leq 1/2$

**Theorem 2.** *Any universal constant  $c$  guaranteeing the existence of an  $\varepsilon$ -light set  $S$  with  $|S| \geq c\varepsilon n$  must satisfy  $c \leq 1/2$ .*

*Proof.* Let  $G$  be a perfect matching on  $n$  vertices (where  $n$  is even). The edge set  $E$  comprises  $n/2$  disjoint edges.

Consider any edge  $e = \{u, v\} \in E$  and the vector  $x = e_u - e_v$ .

- In the full graph  $G$ ,  $u$  and  $v$  have degree 1 and are connected only by  $e$ . The quadratic form is  $x^\top L x = 4$ .
- In the induced subgraph  $G_S$ , if  $\{u, v\} \subseteq S$ , then  $x^\top L_S x = 4$ . Otherwise, if at least one endpoint is missing,  $x^\top L_S x = 0$ .

The condition  $L_S \preceq \varepsilon L$  implies  $x^\top L_S x \leq \varepsilon x^\top L x$ . If  $S$  contains any full edge  $e$ , we have  $4 \leq 4\varepsilon$ , implying  $\varepsilon \geq 1$ .

Thus, for any fixed  $\varepsilon < 1$ , an  $\varepsilon$ -light set  $S$  must be an *independent set*. The maximum size of an independent set in a perfect matching is exactly  $n/2$ . If a universal constant  $c$  existed such that  $|S| \geq c\varepsilon n$ , then:

$$\frac{n}{2} \geq |S| \geq c\varepsilon n \implies c \leq \frac{1}{2\varepsilon}.$$

Since  $\varepsilon$  can be chosen arbitrarily close to 1, we conclude  $c \leq 1/2$ .  $\square$

### 3 Structural Reduction via Linearization

To approach the lower bound, we decouple the quadratic dependence of  $L_S$  on the vertex set  $S$  using a linearization relaxation.

**Lemma 3** (Linearization). *Let  $\delta_u \in \{0, 1\}$  be the indicator for  $u \in S$ . Then:*

$$L_S \preceq \sum_{u \in V} \delta_u B_u, \quad \text{where } B_u := \frac{1}{2} \sum_{v \sim u} L_{\{u, v\}}.$$

*Proof.* We use the inequality  $\delta_u \delta_v \leq \frac{\delta_u + \delta_v}{2}$  for binary variables. Summing over edges:

$$L_S = \sum_{\{u, v\} \in E} \delta_u \delta_v L_{\{u, v\}} \preceq \sum_{\{u, v\} \in E} \frac{\delta_u + \delta_v}{2} L_{\{u, v\}} = \sum_{u \in V} \delta_u B_u.$$

$\square$

#### 3.1 Projection to Image Space

To analyze the spectral norm, we work in the image of  $L$ , denoted  $\text{im}(L)$ . Let  $L^\dagger$  be the Moore-Penrose pseudoinverse. Define the projection  $P = L^{\dagger/2} L L^{\dagger/2}$ , which projects onto  $\text{im}(L)$ . We define the normalized matrices:

$$A_u := L^{\dagger/2} B_u L^{\dagger/2}.$$

These satisfy  $A_u \succeq 0$  and  $\sum_{u \in V} A_u = P$ . All subsequent PSD inequalities are understood to hold on the subspace  $\text{im}(L)$ . Equivalently, one may view the inequalities as holding after applying the projection  $P$  to both sides.

**Theorem 4** (Reduction). *Suppose that for the specific family of matrices  $\{A_u\}$  derived from a graph, and for any  $p \in (0, 1)$ , there exists a subset  $S$  with  $|S| \geq pn$  such that  $\sum_{u \in S} A_u \preceq 2pP$ . Then  $c \geq 1/2$ .*

*Proof.* Set  $p = \varepsilon/2$ . If such a set exists, we have  $|S| \geq (\varepsilon/2)n$  and  $\sum_{u \in S} A_u \preceq \varepsilon P$ .

We now translate this back to the Laplacian scale. Restricted to  $\text{im}(L)$ , we conjugate the inequality by  $L^{1/2}$ :

$$L^{1/2} \left( \sum_{u \in S} A_u \right) L^{1/2} \preceq \varepsilon L^{1/2} P L^{1/2}.$$

Using the fact that  $L^{1/2} L^{\dagger/2}$  acts as the identity on  $\text{im}(L)$ , we have:

$$L^{1/2} A_u L^{1/2} = L^{1/2} (L^{\dagger/2} B_u L^{\dagger/2}) L^{1/2} = B_u.$$

Also,  $L^{1/2} P L^{1/2} = L$ . Thus, the inequality becomes  $\sum_{u \in S} B_u \preceq \varepsilon L$ . By Lemma 3,  $L_S \preceq \sum_{u \in S} B_u \preceq \varepsilon L$ . Thus  $S$  is  $\varepsilon$ -light.

If this hypothesis holds, then  $c \geq 1/2$ . Combined with the upper bound, we would establish  $c = 1/2$ .  $\square$

## 4 The Obstacle and Conjecture

The reduction theorem provides a pathway to  $c = 1/2$ , but the required matrix subset selection hypothesis is **false for general PSD families**.

**Example 5** (Counterexample for General Matrices). *Let  $P = I_d$  and  $A_i = e_i e_i^\top$  for  $i = 1, \dots, d$  (standard basis projectors). Then  $\sum_{i=1}^d A_i = I_d$ . For any non-empty subset  $S$ , the sum  $\sum_{i \in S} A_i$  is a diagonal projection matrix with spectral norm  $\lambda_{\max} = 1$ . The hypothesis requires  $\sum_{i \in S} A_i \preceq 2pI_d$ , implying  $1 \leq 2p$ . This fails whenever  $p < 1/2$ .*

The counterexample relies on  $\{A_i\}$  being mutually orthogonal rank-1 projectors. In the graph case,  $B_u = \frac{1}{2} \sum_{v \sim u} L_{\{u,v\}}$  corresponds to a "half-star" at  $u$ . Crucially, adjacent vertices  $u$  and  $v$  share the edge term  $L_{\{u,v\}}$ , meaning  $B_u B_v \neq 0$ . This structural overlap implies that the normalized matrices  $\{A_u\}$  are *not* mutually orthogonal.

We formalize the missing link as follows:

**Conjecture 6** (Graph-Structured Subset Selection). *Let  $G = (V, E)$  be a graph and let  $\{A_u\}_{u \in V}$  be the normalized half-star matrices defined above. For any  $p \in (0, 1)$ , there exists a subset  $S \subseteq V$  with  $|S| \geq pn$  such that:*

$$\sum_{u \in S} A_u \preceq 2pP.$$

Establishing this conjecture is necessary and sufficient to prove that the sharp universal constant is  $c = 1/2$ .